
Clouding Computing Final Presentation

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Presentation Outline

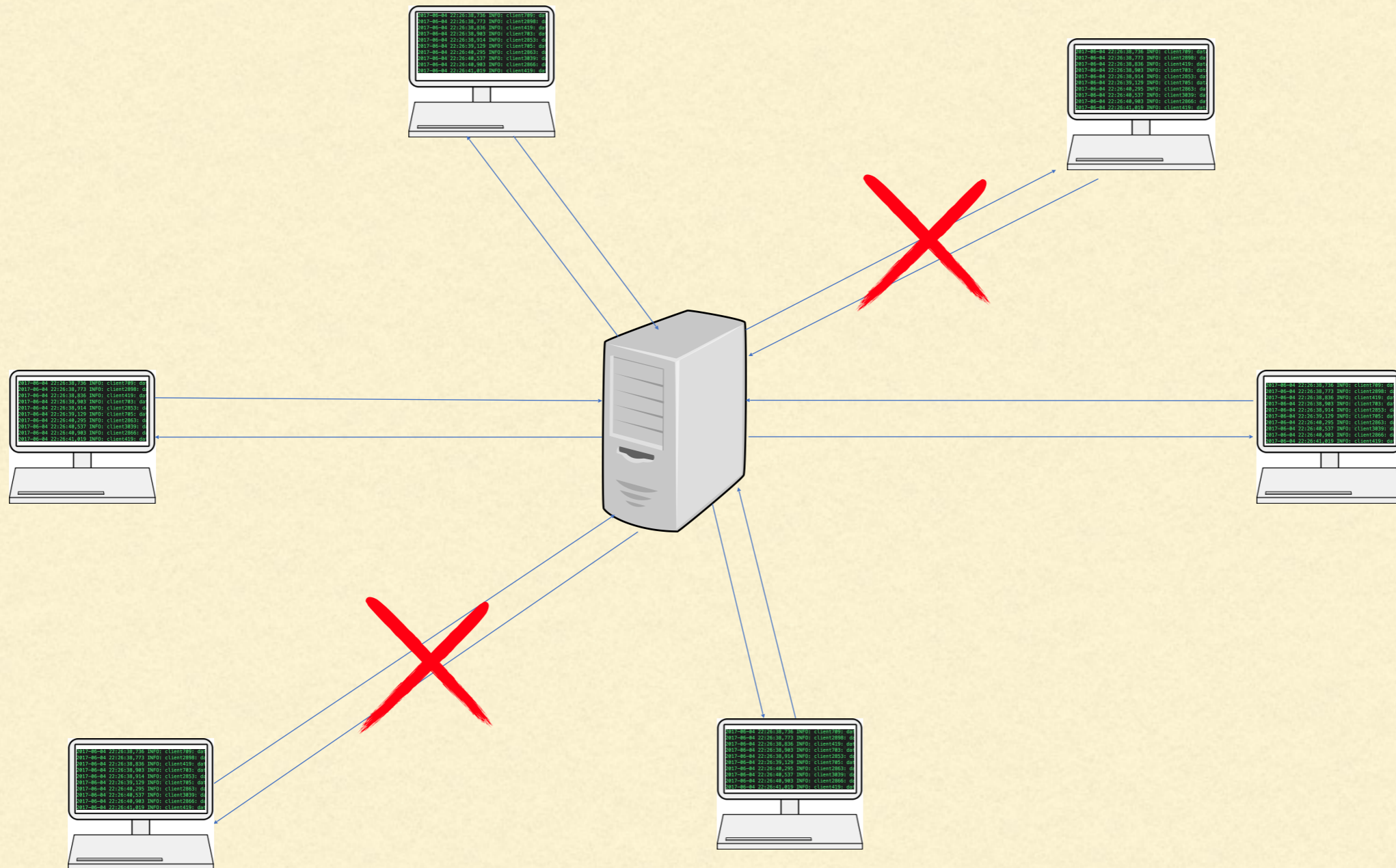
Computation Model

What doesn't work

What does work

System Design and Others

Computation Model



What doesn't work — things that we tried but not included in final implementation

Simulated Annealing

Stochastic Tunneling

Detrended Fluctuation Analysis

Clique Counts Approximation

Simulated Annealing/Stochastic Tunneling

$$e^{-\frac{\Delta f}{kT}} \geq \mathbb{R}(0, 1)$$

$$f_{\text{STUN}}(\vec{x}) = 1.0 - e^{-\frac{\mathbb{E}(\vec{x}) - \mathbb{E}(\vec{x}_0)}{\gamma}}$$

Challenges

- Reset parameters
 - Too slow!!!!
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Detrended Fluctuation Analysis (DFA)

Improving FPGA Placement with
Dynamically Adaptive Stochastic Tunneling

Clique Counts Approximation

A Fast and Provable Method for Estimating
Clique Counts Using Turán's Theorem

Things that works

Tabu Search*

- Use “Tabu Search” to avoid walking to the same graph twice
 - Tabu Search uses linked-list to store history
 - However, we use a **global Bloom Filter** to replace linked-list
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Global Bloom Filter (GBF)

- Client ask GBF whether graph G has been checked before
 - Case 1:
 - If G is indeed checked before, GBF will tell us G has been reached with probability 1
 - Case 2:
 - If G hasn't been checked before, GBF might give a False Positive response (yeah, G has been checked) with probability p
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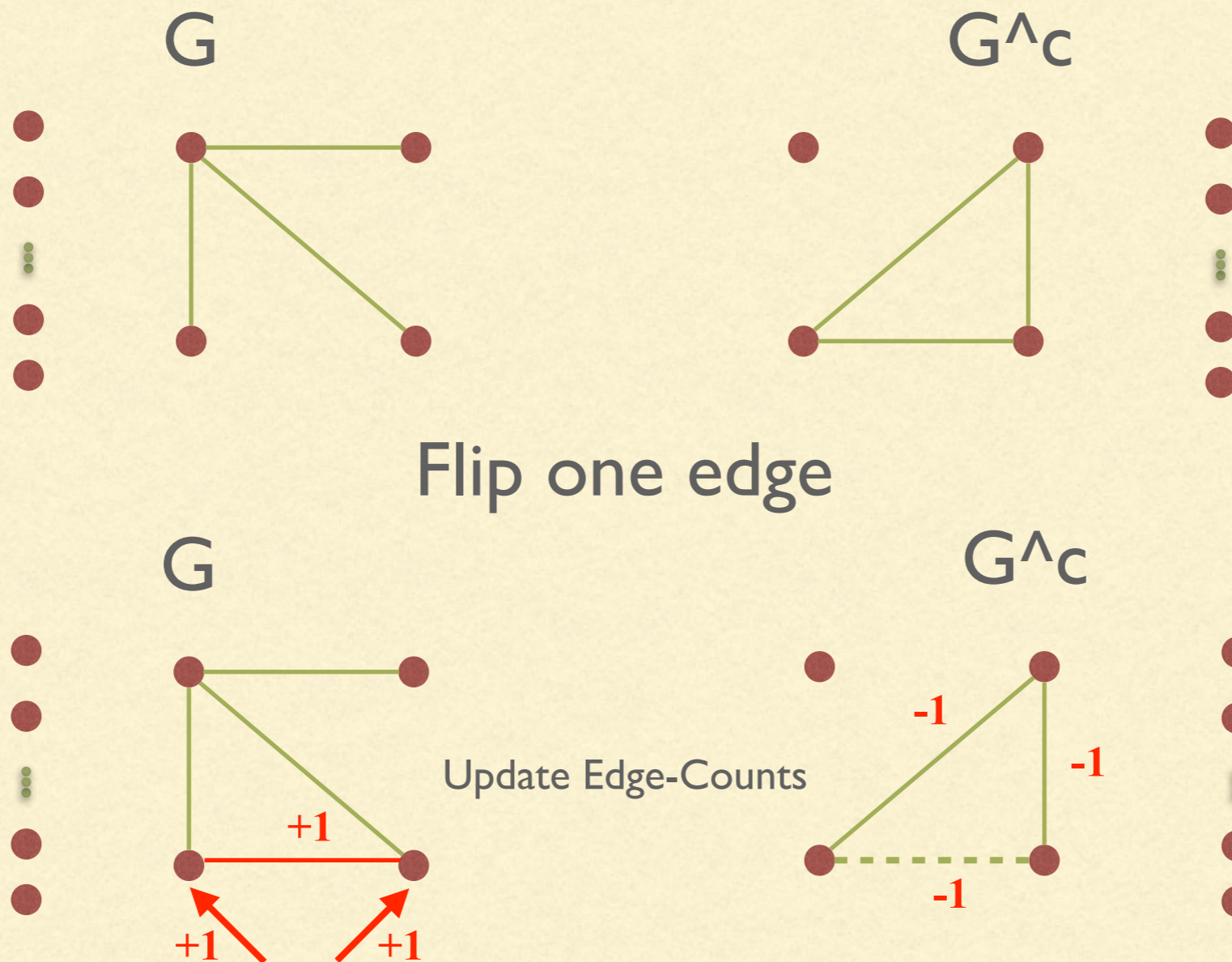
Global Bloom Filter (GBF)

- Benefits of GBF
 - Space efficient because:
 - only store bits (similar to hash) instead of the giant graph/matrix representation
 - Parametrizable because:
 - Configure possible solution number N and false positive rate P
 - For example,
 - $N = 40,000,000, p = 1.0E-10: \text{size} = 228.53\text{MB}$
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The main trick

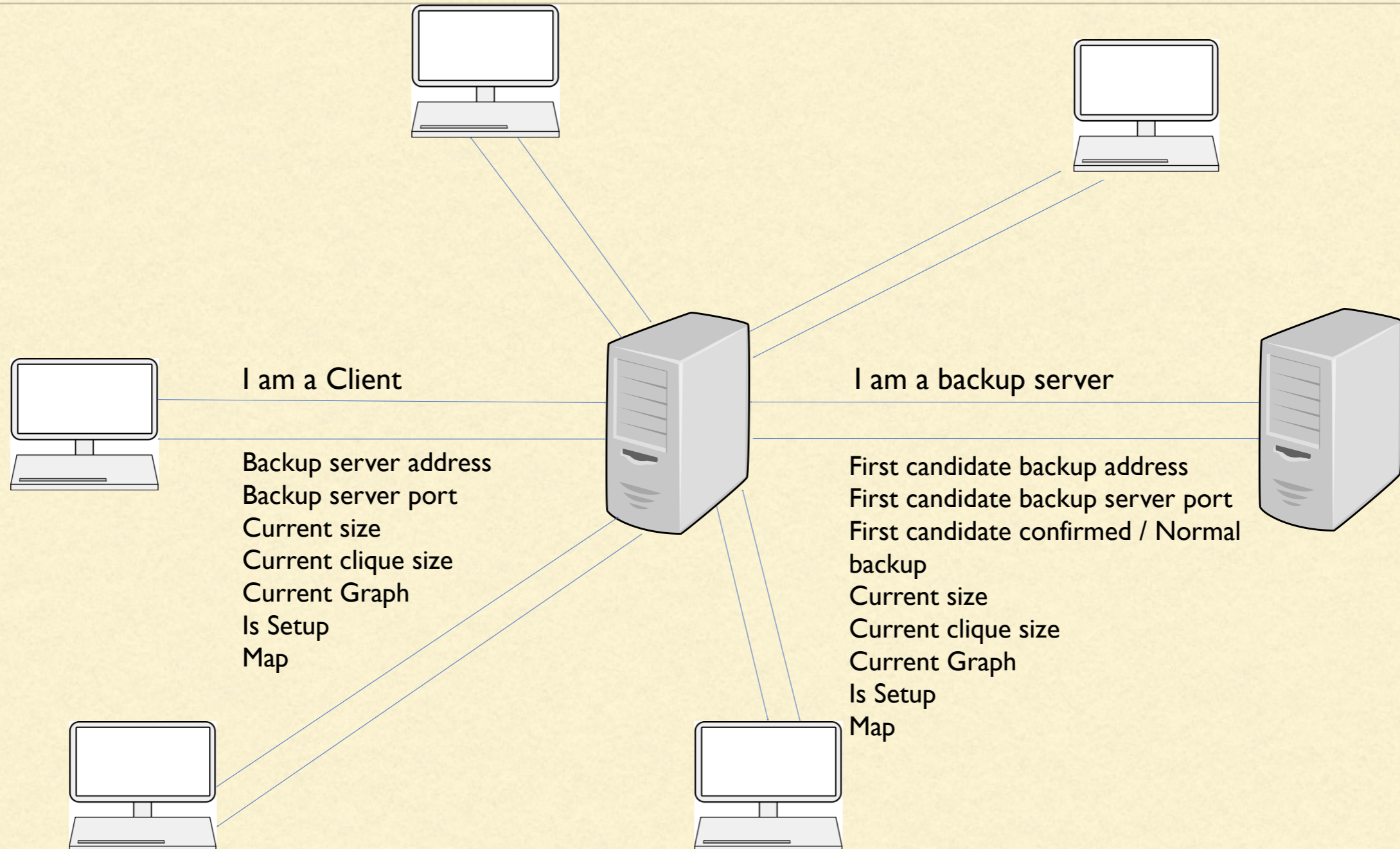
- When move to larger problem size
 - Count 10 cliques once
 - Maintain a Map: {Edges} \rightarrow {clique counts}
 - After the first and only 10-clique count check, we start to flip edges
 - During flipping, We also maintain the “complement” of G , because:
 - To check $R(m,n)$ is the same to check:
 1. m -clique on G
 2. n -clique on G^c
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G's complement and (n-2)-clique

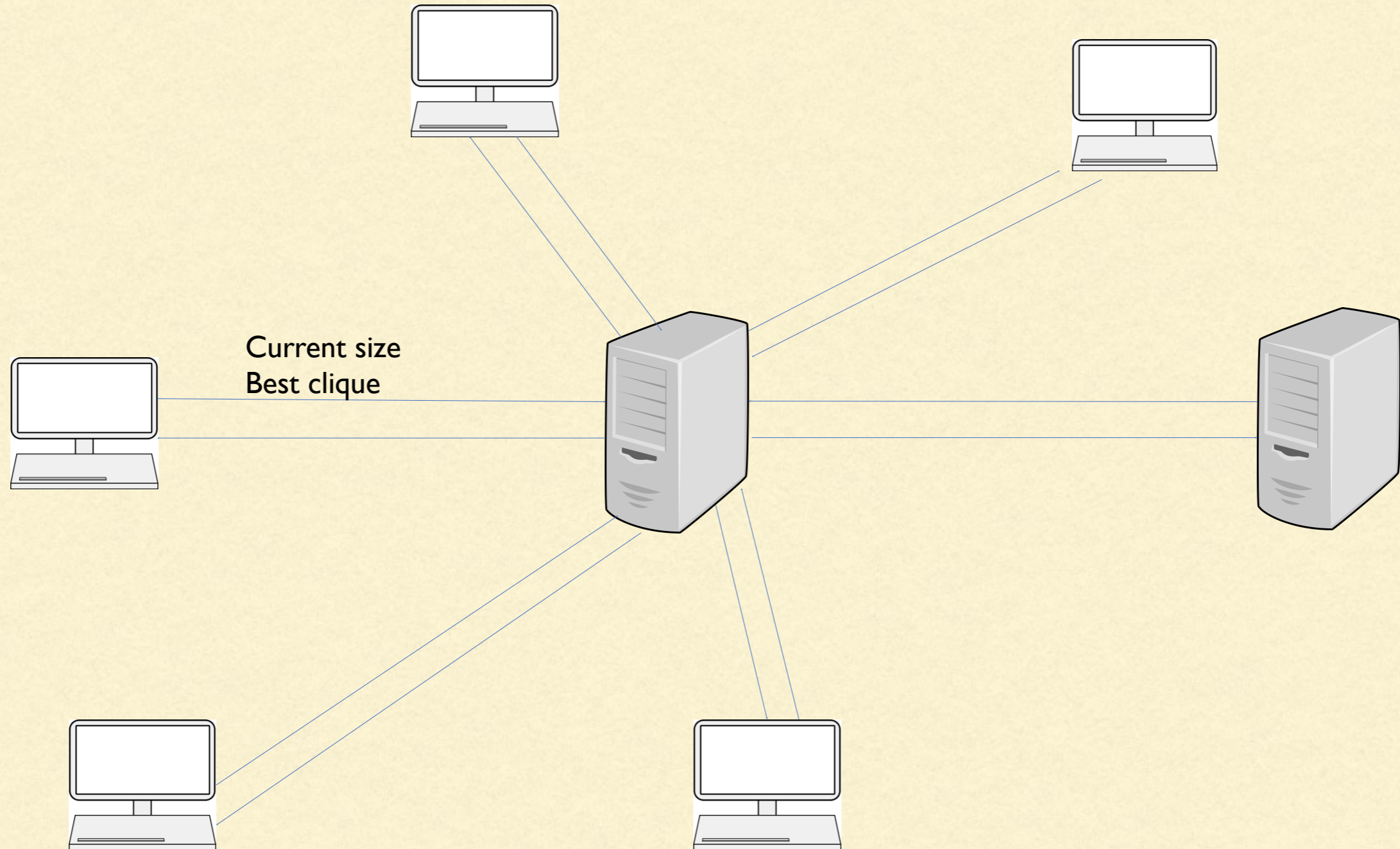


Search for $(n-2)$ -clique in the set of intersected neighbors

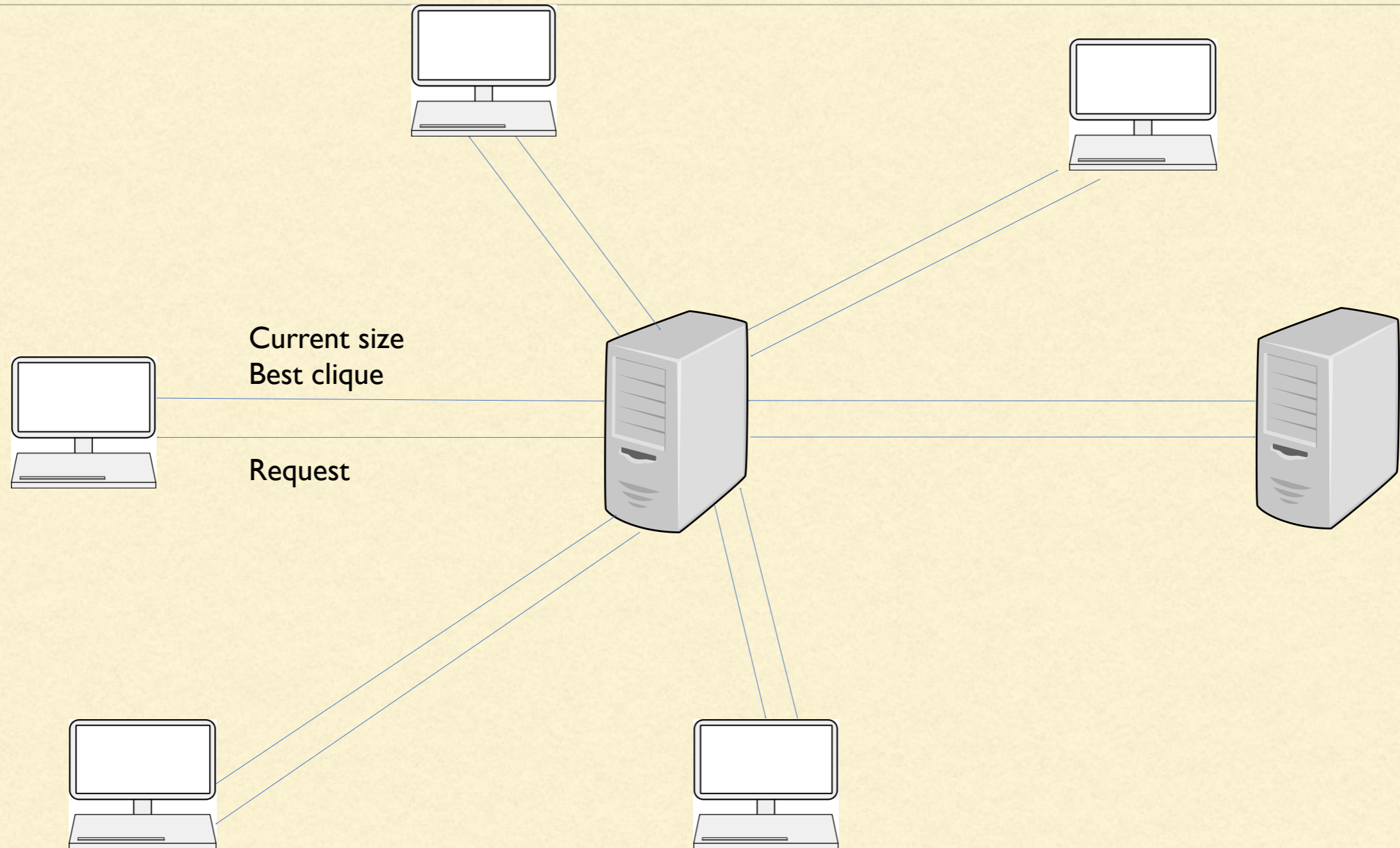
Server-Client Protocol



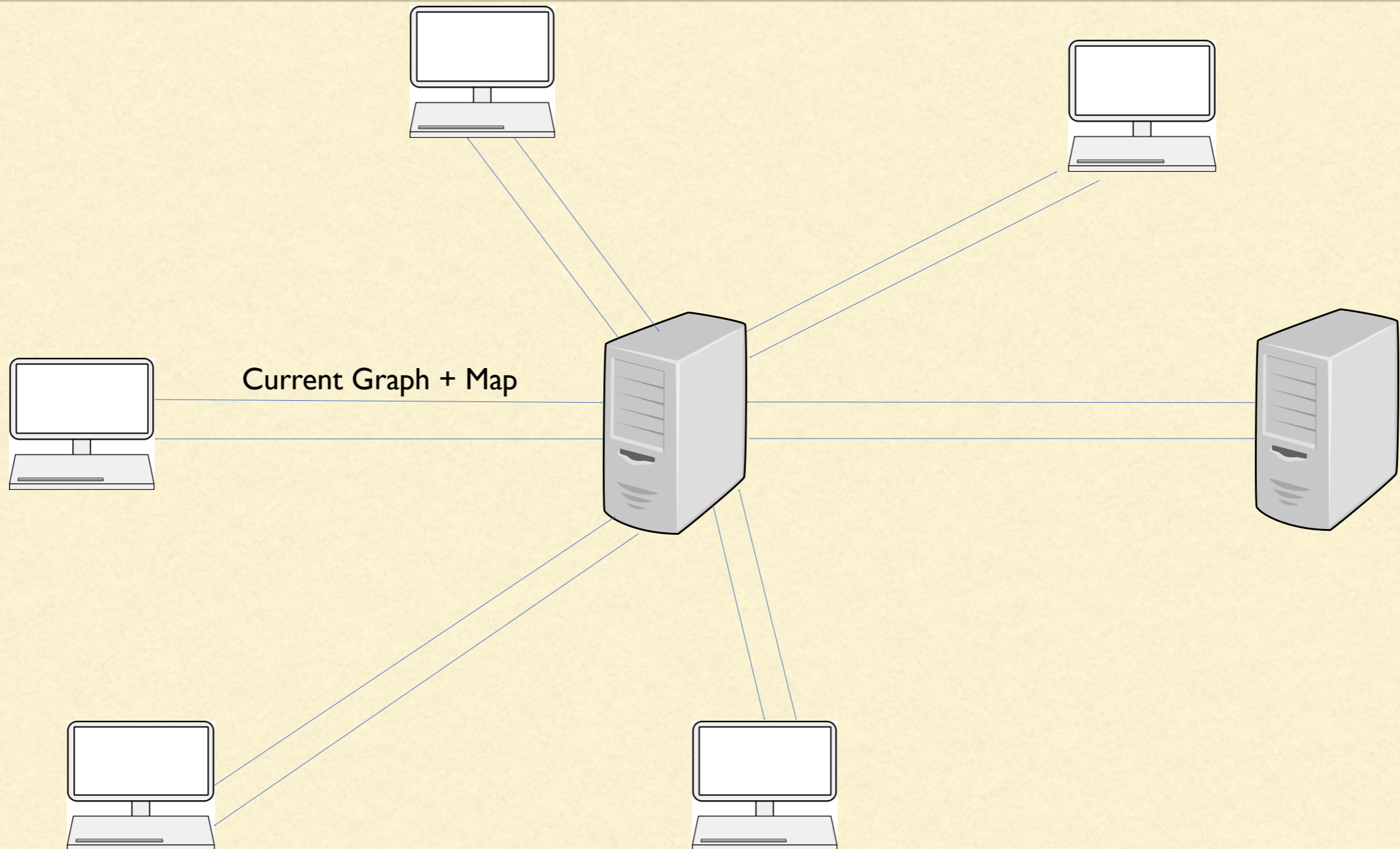
Server-Client Protocol



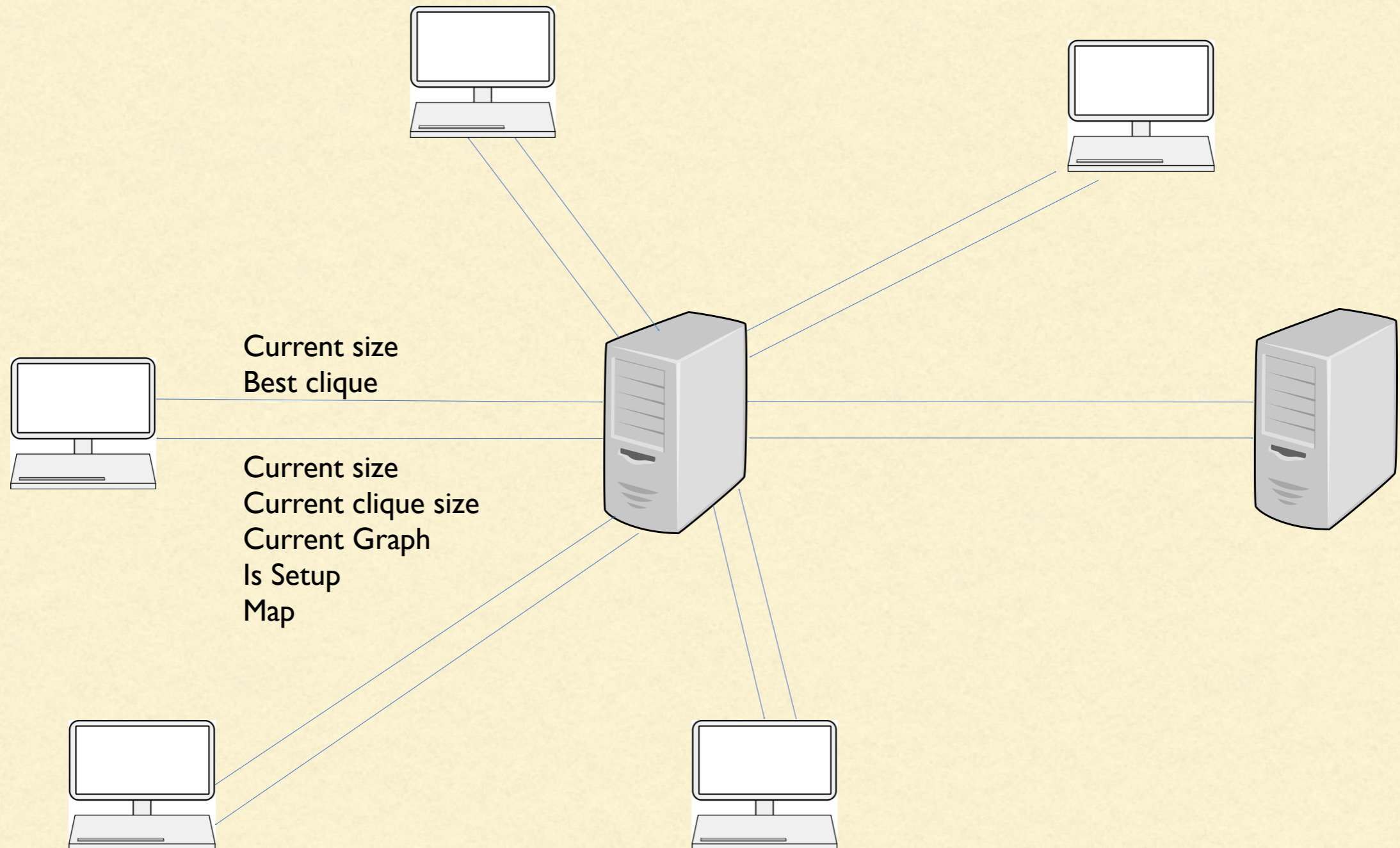
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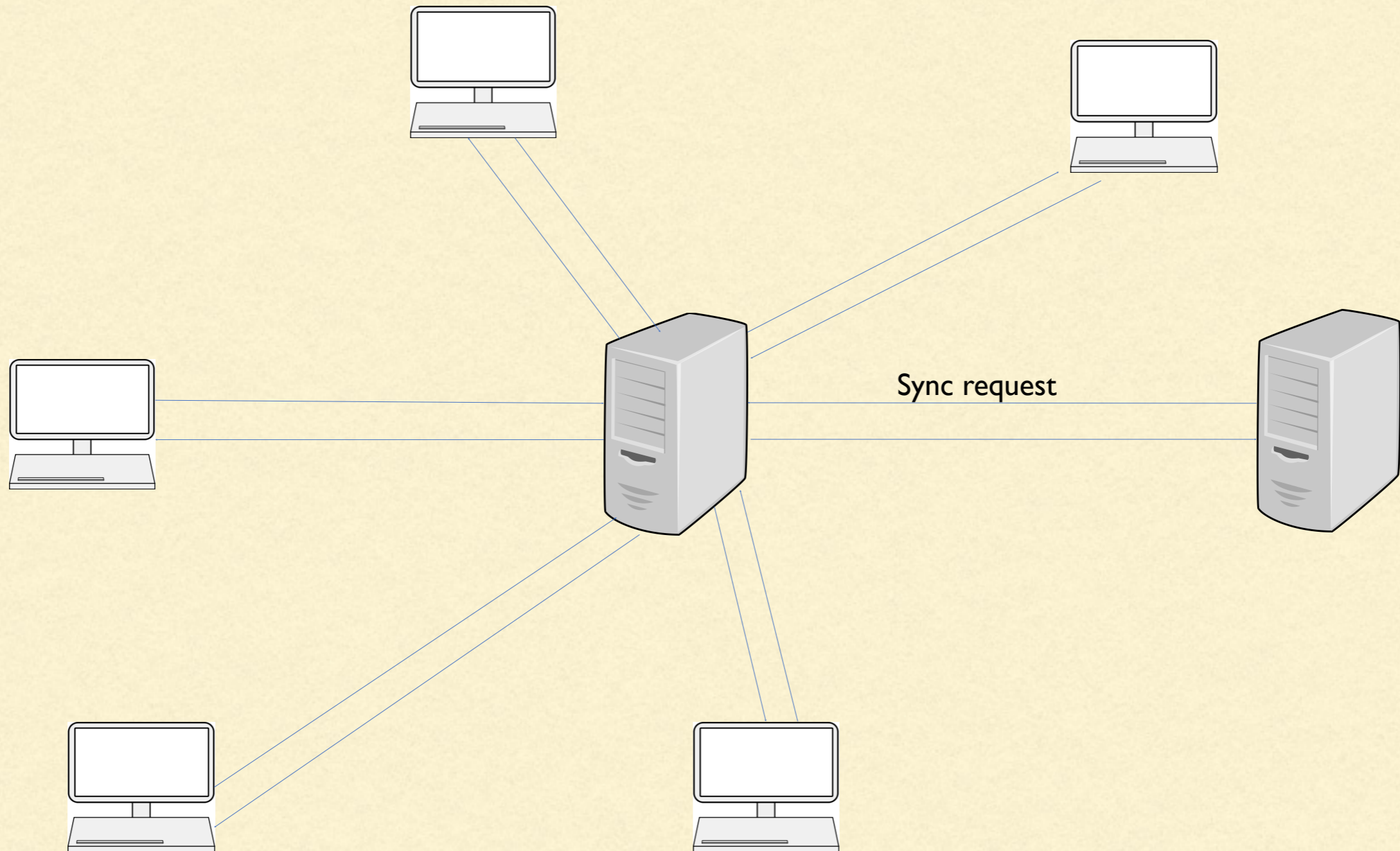
Server-Client Protocol



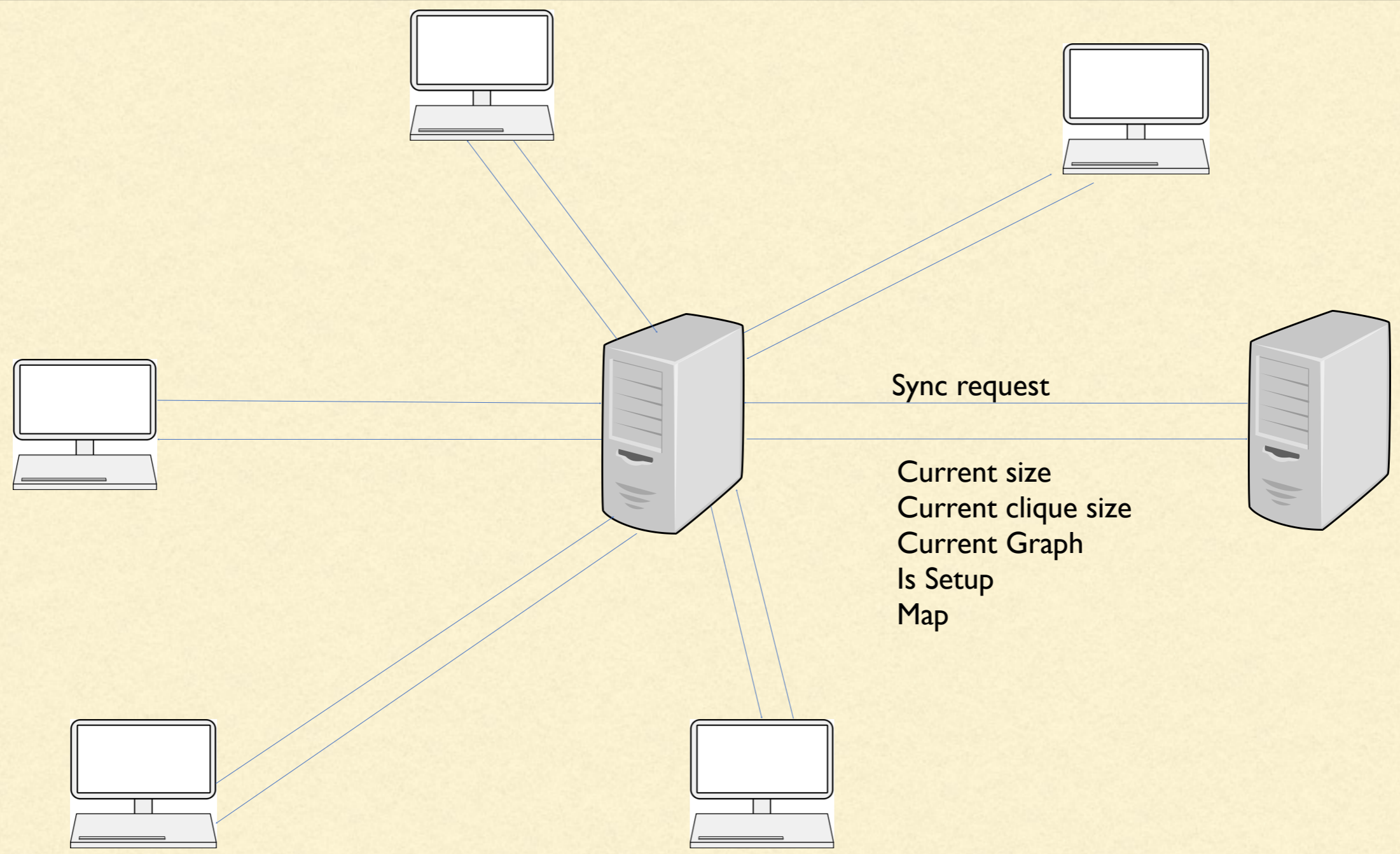
Server-Client Protocol



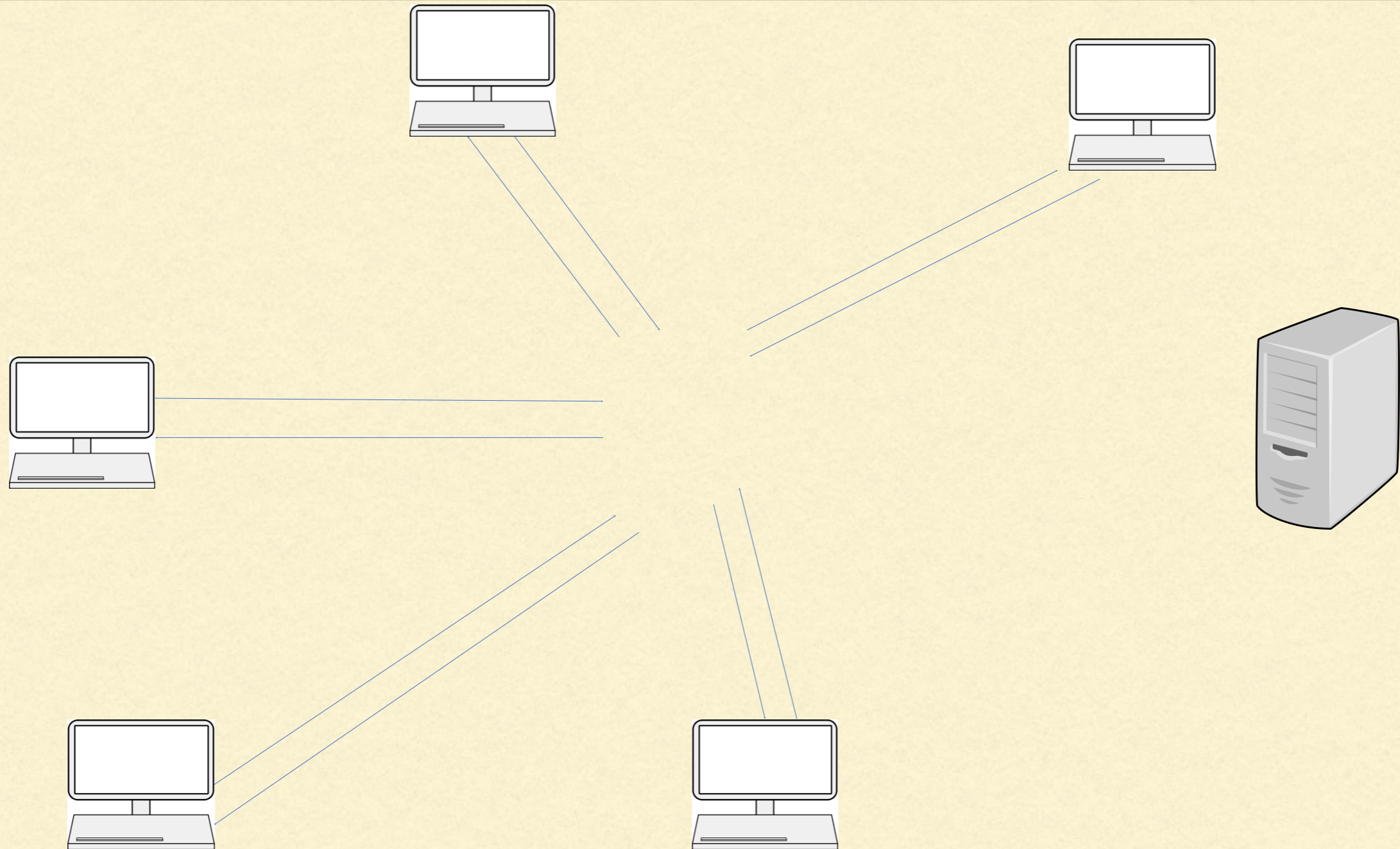
Server-Client Protocol



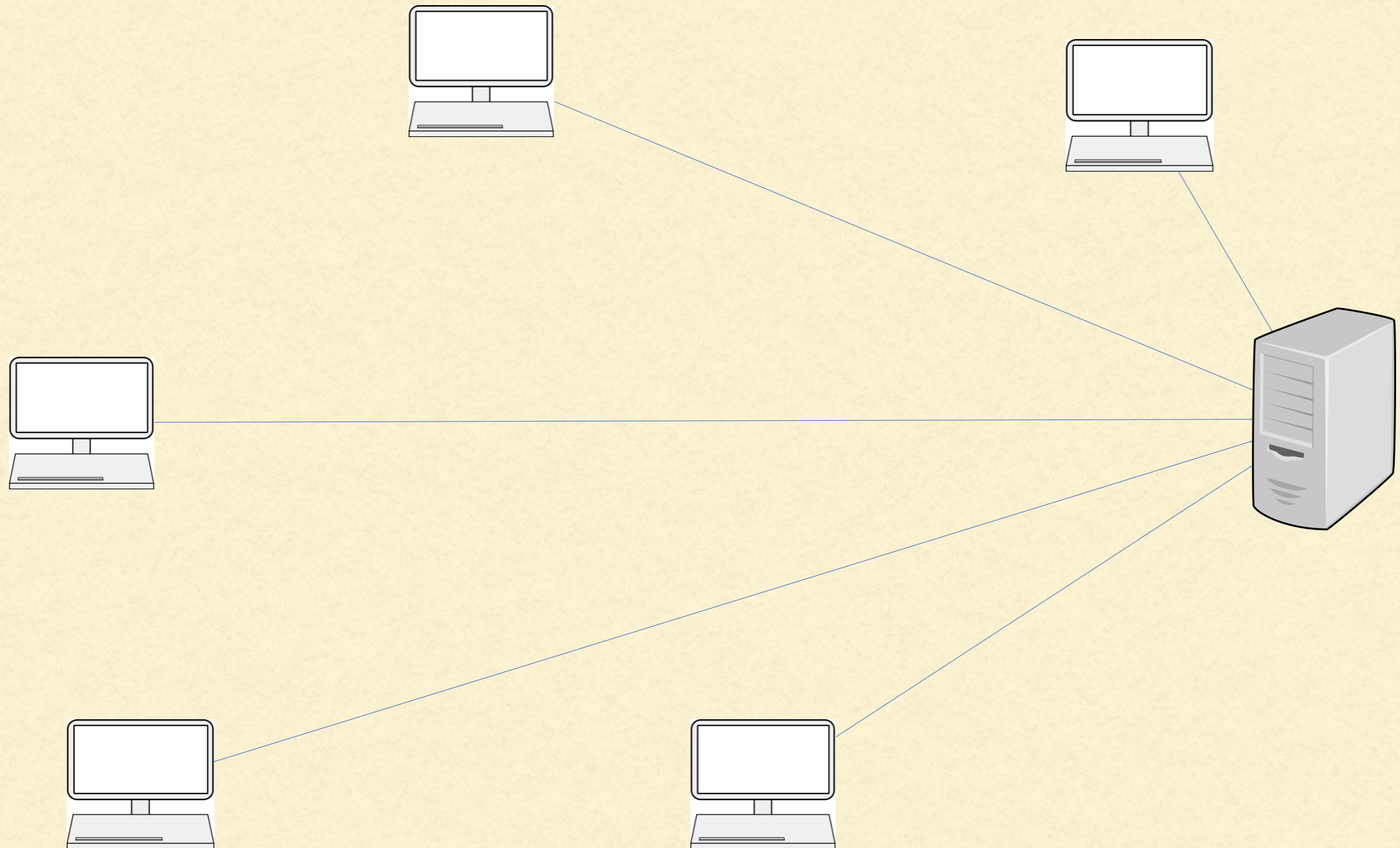
Server-Client Protocol



Server-Client Protocol



Server-Client Protocol



Results

Best Clique Size:
340

CPU counts:

$885842(\text{iteration per day}) * 7 * 10^6(\text{CPU per iteration}) * 35(\text{days})$

$= 2.14 * 10^{14}$

Takeaways

Start early and do more logging
